

# Lecture 7

## Transportation Problem (TP)

special case of Integer Linear  
Programming

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# What is the Transportation Problem?

The transportation problem is a special type of linear programming problem where the objective is to minimize the cost of distributing a product from a number of sources or origins to a number of destinations.

# Transportation Example

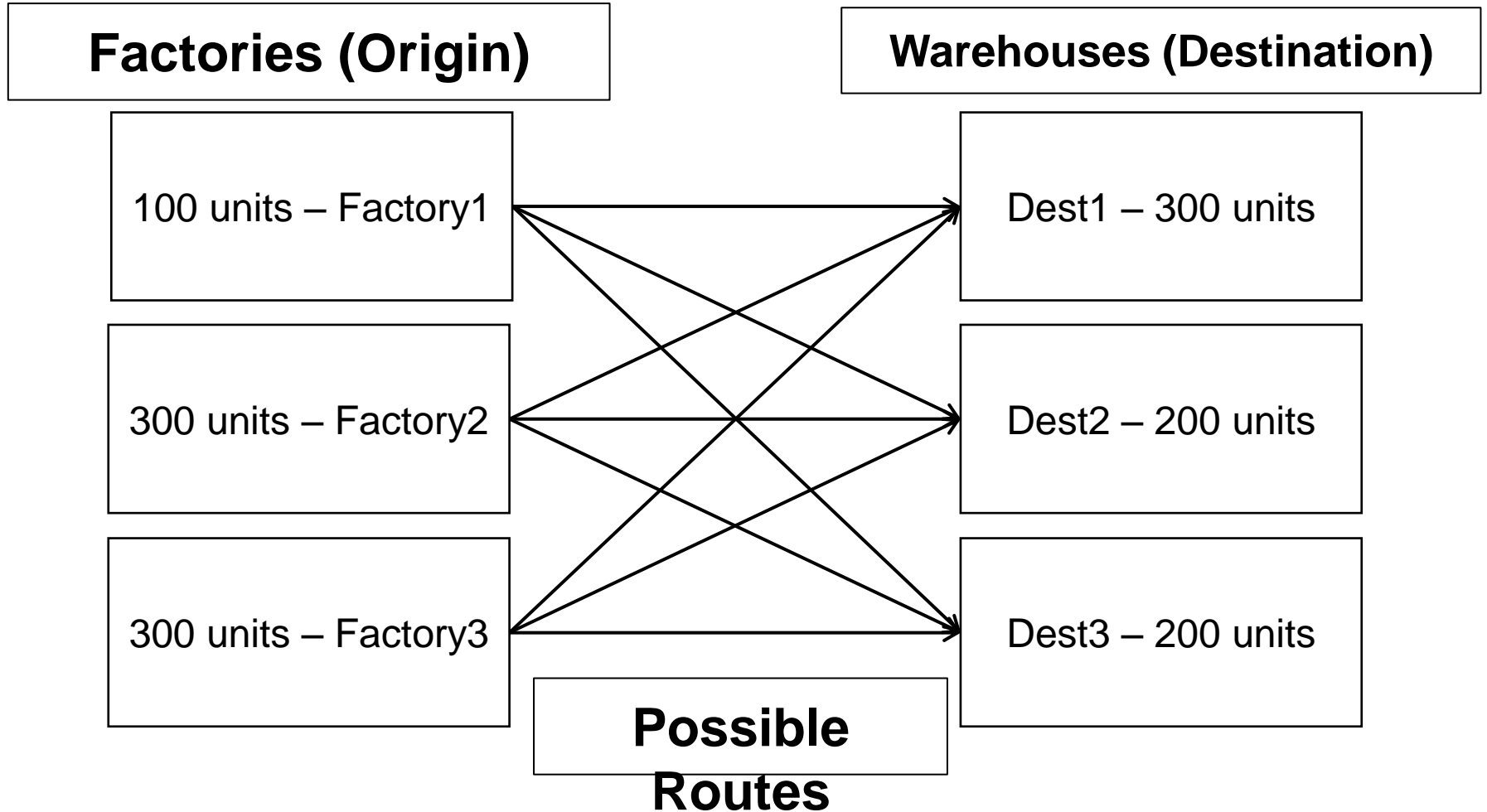
We have 3 factories and 3 warehouses

**Decision:** How much to ship from each origin to each destination?

**Objective:** Minimize transportation cost.

<b>Transportation Cost per Unit</b>			
<u>Factory</u>	<u>Dest 1</u> (300 units)	<u>Dest2</u> (200 units)	<u>Dest3</u> (200 units)
Factory1 (100 units)	Rs. 5	Rs. 4	Rs. 3
Factory2 (300 units)	Rs. 8	Rs. 4	Rs. 3
Factory3 (300 units)	Rs. 9	Rs. 7	Rs. 5

# Transportation Example



## Decision Variables:

$x_{ij}$  = number of units transported from factory “i” to Destination “j”

## Objective Function:

$$\begin{aligned} \text{Min: } Z = & 5x_{11} + 4x_{12} + 3x_{13} \\ & + 8x_{21} + 4x_{22} + 3x_{23} \\ & + 9x_{31} + 7x_{32} + 5x_{33} \end{aligned}$$

## supply Constraints:

$$\begin{aligned} x_{11} + x_{12} + x_{13} & \leq 100 \\ x_{21} + x_{22} + x_{23} & \leq 300 \\ x_{31} + x_{32} + x_{33} & \leq 300 \end{aligned}$$

## demand Constraints:

$$\begin{aligned} x_{11} + x_{21} + x_{31} & \geq 300 \\ x_{12} + x_{22} + x_{32} & \geq 200 \\ x_{13} + x_{23} + x_{33} & \geq 200 \end{aligned}$$

Transportation Cost per Unit			
<u>Factory</u>	<u>Dest1</u> (300 units)	<u>Dest2</u> (200 units)	<u>Dest3</u> (200 units)
Factory1 (100 units)	$x_{11}=?$ $C_{11}=\text{Rs. } 5$	$x_{12}=?$ $C_{12}=\text{Rs. } 4$	$x_{13}=?$ $C_{13}=\text{Rs. } 3$
Factory2 (300 units)	Rs. 8	Rs. 4	Rs. 3
Factory3 (300 units)	Rs. 9	Rs. 7	Rs. 5

# The Transportation Algorithm

# Unbalanced Problems

- If **(Total Supply) > (Total Demand)**,
  - then for Excess supply is assumed to go to the inventory and costs nothing for shipping.
  - Dummy destination column is added, whose demand equals the difference between the total supply and total demand and zero transportation cost
- If **(Total Supply) < (Total Demand)**,
  - *a dummy source* is created, whose supply equals the difference.
  - All unit shipping costs into a dummy destination or out of a dummy source are 0.



# Example 1: Balanced Problem

		DESTINATIONS				Supply
		D <sub>1</sub>	D <sub>2</sub>	D <sub>3</sub>	D <sub>4</sub>	
Sources	S <sub>1</sub>	50	75	30	45	12
	S <sub>2</sub>	65	80	40	60	17
	S <sub>3</sub>	40	70	50	55	11
Demand		10	10	10	10	

## Example 2: Unbalanced Problem

	Destination				Supply
	D1	D2	D3	D4	
S1	50	75	35	75	12
S2	65	80	60	65	17
S3	40	70	45	55	11
(Dummy)	<b>0</b>	<b>0</b>	<b>0</b>	<b>0</b>	<b>10</b>
Demand	15	10	15	10	

# Transportation Tableau:

	D <sub>1</sub>	D <sub>2</sub>	D <sub>3</sub>	D <sub>4</sub>	Supply
S <sub>1</sub>	50	75	30	45	12
S <sub>2</sub>	65	80	40	60	17
S <sub>3</sub>	40	70	50	55	11
Demand	10	10	10	10	

Shipment will be placed here

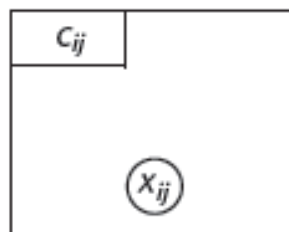
Unit Shipping Cost

■ **TABLE 9.15** Format of a transportation simplex tableau

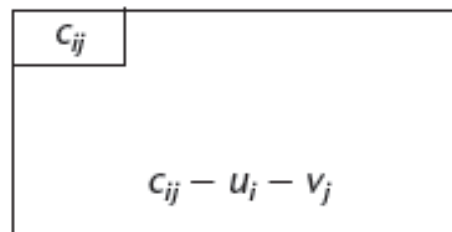
		Destination				Supply	$u_i$
		1	2	...	$n$		
Source	1	$c_{11}$ $x_1$	$c_{12}$ $x_1$	...	$c_{1n}$ $x_1$	$s_1$	
	2	$c_{21}$ $x_2$	$c_{22}$ $x_2$	...	$c_{2n}$	$s_2$	
	⋮	...	...	...	...	⋮	
	$m$	$c_{m1}$ $x_n$	$c_{m2}$	...	$c_{mn}$	$s_m$	
Demand		$d_1$	$d_2$	...	$d_n$	$Z =$	
	$v_j$						

Additional information to be added to each cell:

*If  $x_{ij}$  is a basic variable*



*If  $x_{ij}$  is a nonbasic variable*



# The Transportation Problem

- Transportation method is more efficient
  - Especially for large problems
- For transportation problems with  $m$  sources and  $n$  destinations:
  - Number of basic variables is equal to  $m + n - 1$

# The Transportation Method

The method involves **3 steps**:

## **1. Obtaining Initial Basic Feasible Solution**

- a. North-West Corner Rule
- b. Vogel's Approximation Method

## **2. Testing the Optimality**

## **3. Improving the Solution**

# 1- Initial Solution Procedure:

## 1. Northwest Corner Starting Procedure

1. Select the variable in the upper left (northwest) corner
2. Allocate the minimum of  $s$  or  $d$  to this variable. If this minimum is  $s$ , eliminate all variables in its row from future consideration and reduce the demand in its column by  $s$ ; if the minimum is  $d$ , eliminate all variables in the column from future consideration and reduce the supply in its row by  $d$ .

REPEAT THESE STEPS UNTIL ALL SUPPLIES HAVE BEEN ALLOCATED.

# Example 1:

	D <sub>1</sub>	D <sub>2</sub>	D <sub>3</sub>	D <sub>4</sub>	Supply
S <sub>1</sub>	50	75	<del>30</del>	<del>45</del>	<del>12</del> 2 0
S <sub>2</sub>	<del>65</del>	80	40	<del>60</del>	<del>17</del> 9 0
S <sub>3</sub>	<del>40</del>	<del>70</del>	50	<del>55</del>	<del>11</del> 10 0
Demand	<del>10</del>	<del>10</del>	<del>10</del>	<del>10</del>	
	0	8	1	0	
		0	0		

A 3x4 transportation problem matrix with supply and demand values. The matrix is partially filled with blue cells and contains red arrows indicating adjustments. The supply values are 12, 17, and 11, and the demand values are 10, 10, 10, and 10. The total shipping cost is 2250.

Total shipping cost = 2250



**Example 2:** Solve the Transportation Table to find Initial Basic Feasible Solution using North-West Corner Method.

$$\begin{aligned} \text{Total Cost} &= 19 \cdot 5 + 30 \cdot 2 + 30 \cdot 6 + 40 \cdot 3 + 70 \cdot 4 + 20 \cdot 14 \\ &= \text{Rs. } 1015 \end{aligned}$$

	D1	D2	D3	D4	Supply
S <sub>1</sub>	19	30	50	10	7
	5	2			
S <sub>2</sub>	70	30	40	60	9
		6	3		
S <sub>3</sub>	40	8	70	20	18
			4	14	
Demand	5	8	7	14	34

# Limitations of North West Corner Rule

- Although this method is relatively **simple**, it is not efficient in terms of cost minimizing. Because it takes into account only the available supply and demand requirements in making assignments and takes no account of the transportation cost involved.

# The Transportation Method

Broadly the method involves 3 steps:

## 1. Obtaining Initial Basic Feasible Solution

- a. North-West Corner Rule ✓
- b. Vogel's Approximation Method

## 2. Testing the Optimality

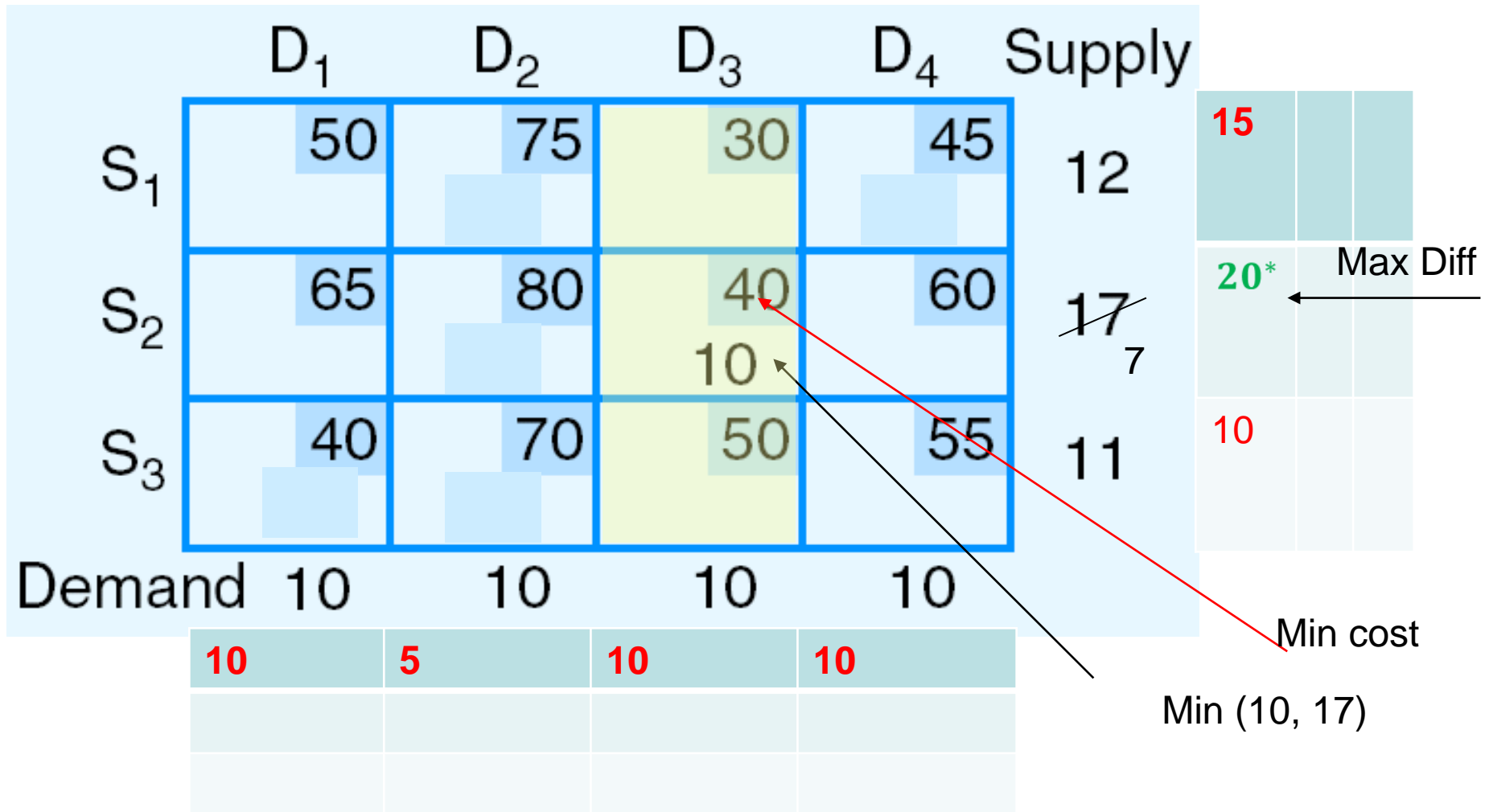
## 3. Improving the Solution

### 3. Vogel's Approximation Method Starting Procedure

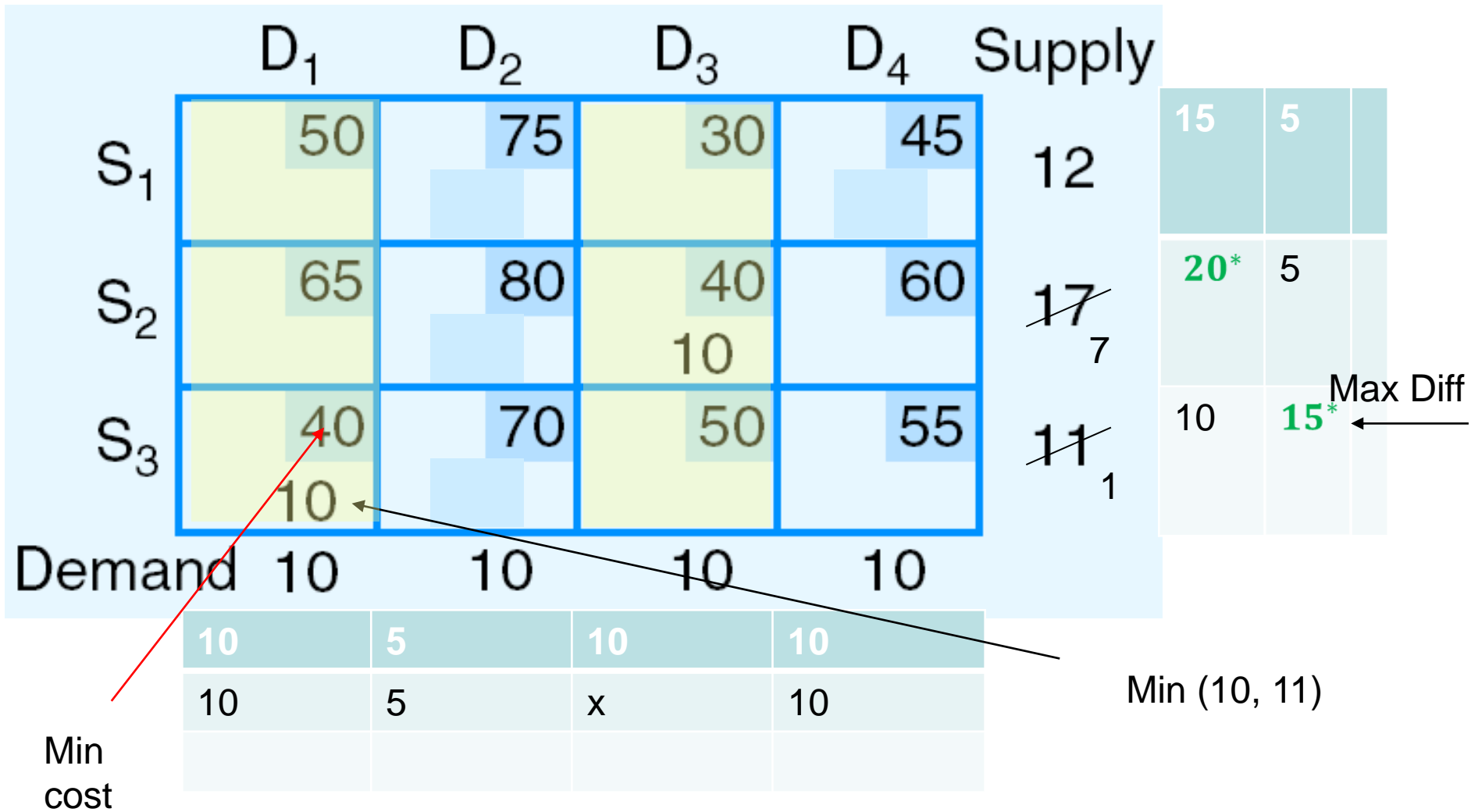
1. For each remaining row and column, determine the **difference between the lowest two remaining costs**; these are called the *row and column penalties*.
2. Select the row or column with the **largest difference** found in step 1 and note the supply remaining for its row,  $s$ , and the demand remaining in its column,  $d$ .
3. Allocate the minimum of  **$s$**  or  **$d$**  to the variable in the selected row or column with the lowest remaining unit cost. If this minimum is  $s$ , eliminate all variables in its row from future consideration and reduce the demand in its column by  $s$ ; if the minimum is  $d$ , eliminate all variables in the column from future consideration and reduce the supply in its row by  $d$ .

REPEAT THESE STEPS UNTIL ALL SUPPLIES HAVE BEEN ALLOCATED.

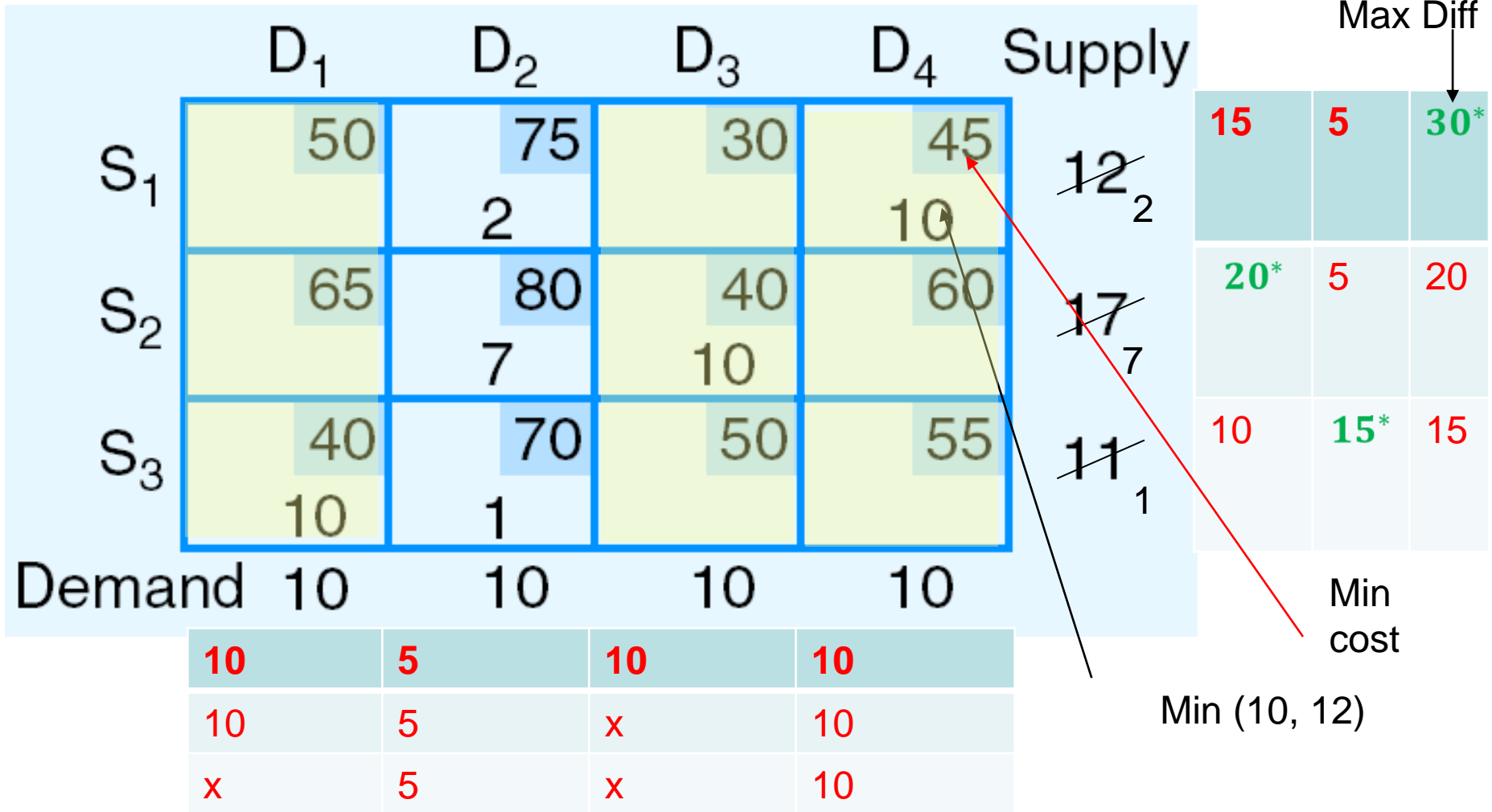
# Example 1



# Example 1 cont.



# Example 1 cont.



Total shipping cost = 2030 **while** in NW Total shipping cost = 2250

# Example 2:

	D1		D2		D3		D4		Supply	Row Diff.
S1	19		30		50		10		7	9
S2	70		30		40		60		9	10
S3	40		8		70		20		<del>18</del> 10	12
Demand	5		8		7		14		34	
Col. Diff.	<b>21</b>		<b>22</b>		<b>10</b>		<b>10</b>			

1- Max difference

2- Min cost

3- Max allocation=  
 $\text{Min}(8, 18)$



## Example 2:

	D1	D3	D4	Supply	Row Diff.
S1	19	50	10	7	<b>9</b>
	5				
S2	70	40	60	9	<b>20</b>
S3	40	70	20	10	<b>20</b>
Demand	5	7	14	34	
Col. Diff.	<b>21</b>	<b>10</b>	<b>10</b>		

## Example 2:

	D3		D4		Supply	Row Diff.
S1	50		10		2	<b>40</b>
S2	40		60		9	<b>20</b>
S3	70		20		10	<b>50</b>
				<b>10</b>		
Demand	7		14		34	
Col.Diff.	<b>10</b>		<b>10</b>			

## Example 2:

	D3	D4	Supply	Row Diff.
S1	50	10	2	<b>40</b>
S2	40	60	9	<b>20</b>
Demand	7	4	34	
Col. Diff.	<b>10</b>	<b>50</b>		

## Example 2:

	D3	D4	Supply	Row Diff.
S2	40	60	9	<b>20</b>
	7	2		
Demand	7	2	34	
Col. Diff.				



## Example 2:

The total transportation cost obtained by this method  
=  $8*8+19*5+20*10+10*2+40*7+60*2$   
= Rs.779

while by using Northwest



Total Cost =  $19*5+30*2+30*6+40*3+70*4+20*14$   
= Rs. 1015

Here, we can see that ***Vogel's Approximation Method*** involves the lowest cost than *North-West Corner Method* hence is the most preferred method of finding initial basic feasible solution.

# The Transportation Method

Broadly the method involves 3 steps:

## 1. Obtaining Initial Basic Feasible Solution

- a. North-West Corner Rule 
- b. Vogel's Approximation Method 

## 2. Testing the Optimality

## 3. Improving the Solution

## 2. Testing the Optimality

Find an initial basic feasible solution by some starting procedure. Then,

1. Set  $U_i = 0$  or  $V_j = 0$  the row or column with the largest number of basic variables. Solve for the other  $U_i$ 's and  $V_j$ 's by:

$$C_{ij} = U_i + V_j \quad \text{for basic variables.}$$

Then calculate the  $C_{ij}-Z_{ij}$  values for non-basic variables by:

$$C_{ij} - Z_{ij} = C_{ij} - U_i - V_j$$

A Basic Feasible Solution is Optimal if  $C_{ij} - U_i - V_j \geq 0$  for each  $(i, j)$

- If all  $C_{ij}-Z_{ij}$  values are nonnegative, STOP; the current solution is optimal.

Else, the current solution is not optimal then,



## Improve the solution

- 1- Choose the non-basic variable with the most negative  $C_{ij} - U_i - V_j$  value as the entering variable
2. Find the cycle that includes the entering variable and some of the BASIC variables.

### Cycle Properties:

- 1- Begin with the Most negative  $C_{ij} - U_i - V_j$  (Entering Variable)
  - 2- All the elements of the loop are basic Variable except the beginning one (E.V.).
  - 3- The number of the elements in the Loop even (min 4 elements)
  - 4- There is no consecutive 3 elements in the same row or the same column.
  - 5- move to left or right or top or bottom but not Diagonal
- > There is a Unique Loop in each basic feasible Solution

## **Note:**

**There must be  $m + n - 1$  basic variables for the transportation simplex method to work!**

**=> Add dummy source or dummy destination, if necessary**

**( $m$ =# of sources and  $n$ =# of destinations)**

# Example 1:

$C_{ij} - U_i - V_j$

Entering Variable

	D <sub>1</sub>	D <sub>2</sub>	D <sub>3</sub>	D <sub>4</sub>	Supp	U <sub>i</sub>
S <sub>1</sub>	50 +	75 2	30 -5	45 10	12	75
S <sub>2</sub>	65 +	80 7	40 10	60 +	17	80
S <sub>3</sub>	40 10	70 1	50 +	55 +	11	70
Demand	10	10	10	10		
V <sub>j</sub>	-30	0	-40	-30		

Total shipping cost = 2030

# Example 1:

Step 2: Determine Which Current Basic Variable Reaches 0 First

	D <sub>1</sub>	D <sub>2</sub>	D <sub>3</sub>	D <sub>4</sub>	Supply
S <sub>1</sub>	50	-75	+30	45	12
S <sub>2</sub>	65	+80	-40	60	17
S <sub>3</sub>	40	70	50	55	11
Demand	10	10	10	10	

Diagram illustrating the cycle property: A cycle is formed by the cells (S<sub>1</sub>, D<sub>2</sub>), (S<sub>1</sub>, D<sub>3</sub>), (S<sub>2</sub>, D<sub>3</sub>), and (S<sub>2</sub>, D<sub>2</sub>). The values in these cells are 2, 10, 10, and 7 respectively. Blue arrows indicate the cycle: from (S<sub>1</sub>, D<sub>2</sub>) to (S<sub>1</sub>, D<sub>3</sub>), then to (S<sub>2</sub>, D<sub>3</sub>), then to (S<sub>2</sub>, D<sub>2</sub>), and finally back to (S<sub>1</sub>, D<sub>2</sub>).

Note: 1. Cycle property

2. X<sub>12</sub> is the leaving variable

All Recipient + Min Donor  
All Donor – Min Donor

Min Donor=2

		Entering Variable
Donor	Recipient	
2-2	2	
Recipient	Donor	
7+2	10-2	

# Example 1:

## Step 3: Determine the Next Transportation Tableau

	D <sub>1</sub>	D <sub>2</sub>	D <sub>3</sub>	D <sub>4</sub>	Supply	U <sub>i</sub>
S <sub>1</sub>	(+) 50 2	(+) 75 9	30 8	45 10	12	0
S <sub>2</sub>	(+) 65 10	80 1	40 50	(+) 60 55	17	10
S <sub>3</sub>	40 10	70 1	(+) 50 10	(+) 55 10	11	0
Demand	10	10	10	10		
	V <sub>j</sub>	40	70	30	45	

Total shipping cost = 2020

{ Improvement = 2 (-5) = 10 }

$C_{ij} - U_i - V_j \geq 0$  for each  $(i, j) \rightarrow$  **Optimal Soln.**

Departing  
Variable

Donor

X

Entering Variable

Recipient

2

Recipient

Donor

9

8

## Example 1:

$$\begin{aligned}\text{Total shipping cost} &= 30*20+45*10+9*80+8*40+10*40+70*1 \\ &= 2020\end{aligned}$$

*Since  $C_{ij} - U_i - V_j \geq 0$ , for each  $(i, j) \rightarrow$  Optimal Soln*

# Example 2:

	D1		D2		D3		D4		Supply	ui
S1	19		30		50		10		7	$c_{ij}-v_j=10$
		5		(+)		(+)		2		
S2	70		30	E.V	40		60		9	60
		(+)		-18		7		2		
S3	40		8		70		20		18	20
		(+)			8		(+)	10		
Demand		5		8		7		14	34	
$v_j$		9		-12		-20		0		

**Loop:**

Recipient

□

Donor

2

Min Donor

Min Donor will be the Departing Variable

Donor

8

Recipient

10



Recipient

+2

Donor

2-2

Donor

8-2

Recipient

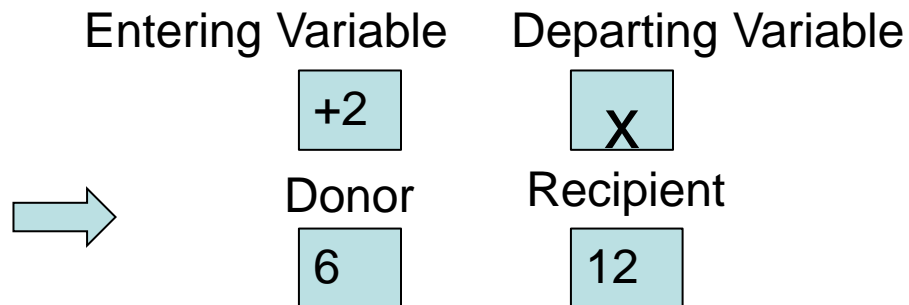
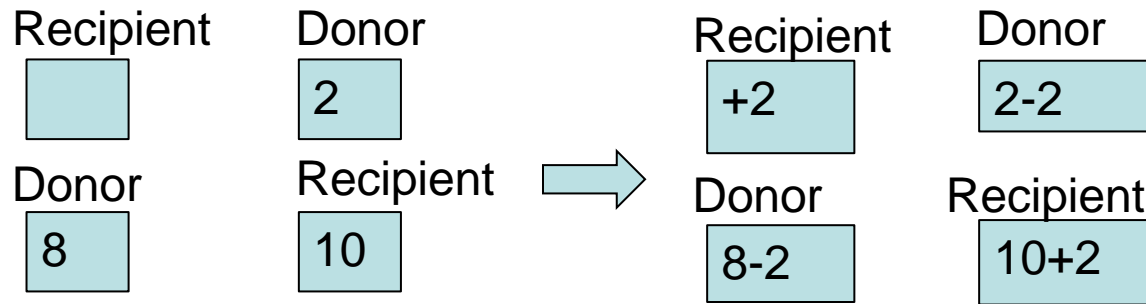
10+2

All Recipient + Min Donor

All Donor - Min Donor

# Example 2:

All Recipient + Min Donor  
All Donor – Min Donor





# Example 2:

Entering Variable      Departing Variable

+2

X

Donor

Recipient

6

12

	D1		D2		D3		D4		Supply	ui
S1	19		30		50		10		7	0
		5		(+)		(+)		2		
S2	70		30	E.V	40		60		9	32
		(+)		2		7		(+)		
S3	40		8		70		20		18	10
		(+)		6		(+)		12		
Demand	5		8		7		14		34	
vj	19		-2		8		10			

## Example 2:

The total transportation cost before improvement  
 $= 8*8+19*5+20*10+10*2+40*7+60*2$   
 $= \text{Rs.779}$

The total transportation cost after improvement  
 $= 19*5+10*2+30*2+40*7+8*6+20*12$   
 $= \text{Rs. 743}$